

Refined Cyclic Sieving on Words and Tableaux

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based on joint work with
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Outline

- ▶ The cyclic sieving phenomenon (CSP) and refinements

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The Cyclic Sieving Phenomenon (CSP)

Definition (Reiner–Stanton–White, 2004)

Take $(X, C, f(q))$ where X is a finite set, C is a finite cyclic group acting on X , and $f(q) \in \mathbb{Z}_{\geq 0}[q]$.

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We say $(X, C, f(q))$ exhibits *the cyclic sieving phenomenon (CSP)* if for all $c \in C$ and roots of unity $\omega \in \mathbb{C}$ of the same order as c ,

$$\#\{x \in X : c \cdot x = x\} = f(\omega).$$

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(Equivalently, $f(\omega)$ is $\text{Tr}_{\mathbb{C}\{X\}}(c)$. Note $f(1) = \#X$.)

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Example

Let $X = \binom{[n]}{k}$ and let $C = \mathbb{Z}/n$ act on X by addition mod n : if $n = 6, k = 3$, then

$$\bar{2} \cdot \{2, 3, 5\} = \{4, 5, 1\}.$$

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Recall:

- ▶ $\binom{(n)}{k}_q := \frac{[n]_q!}{[k]_q! [n-k]_q!}$
- ▶ $[n]_q! := [n]_q [n-1]_q \cdots [1]_q$
- ▶ $[c]_q := 1 + q + \cdots + q^{c-1}$

CSP Refinements

Notation

Given $\text{stat}: X \rightarrow \mathbb{Z}_{\geq 0}$, write

$$\chi^{\text{stat}}(q) := \sum_{x \in X} q^{\text{stat}(x)} \in \mathbb{Z}_{\geq 0}[q].$$

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In *many* CSP triples, $f(q) = X^{\text{stat}}(q)$ for some stat .

Example

$$\binom{n}{k}_q = \binom{[n]}{[k]}^{\text{Sum}'}(q) \text{ where } \text{Sum}'(A) = \left(\sum_{a \in A} a\right) - (1 + 2 + \cdots + k).$$

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Example

Take $X = \binom{[6]}{3}$, $Y = \mathbb{Z}/6 \cdot \{2, 3, 4\}$. Then

$Y^{\text{Sum}'}(q) = 1 + 2q^3 + 2q^6 + q^9$, and

$$Y^{\text{Sum}'}(1) = 6, \quad Y^{\text{Sum}'}(-1) = 0,$$

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We would need $Y^{\text{Sum}'}(\omega_3) = 0$, not 6. So, $(Y, \mathbb{Z}/n, Y^{\text{Sum}'}(q))$ does NOT quite refine the RSW CSP $(X, \mathbb{Z}/n, X^{\text{Sum}'}(q))$.

A First Refinement Result

The *cyclic blocks* of a subset of $[n]$ are maximal sequences of adjacent elements in the subset, where 1 is considered adjacent to n . (Ex: $\{1, 2, 4, 6\} \subset [6]$ has two cyclic blocks, 612 and 4.)

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Here $X^{\text{Sum}'}(q)$ is “equivalent” to $\binom{n}{k}_q$, so the unrefined triple is essentially RSW’s.

Word Combinatorics

Definition

Given a word $w = w_1 \cdots w_n$ with letters $w_i \in \mathbb{Z}_{\geq 1}$, the *descent set* of w is

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(Ex: If $w = 323314$, then $\text{Des}(w) = \{1, 4\}$, $\text{maj}(w) = 1 + 4 = 5$, and $\alpha = (1, 1, 3, 1)$.)

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Remark

They actually proved a generalization valid for all finite Coxeter groups using Springer's regular elements, representation theory, coinvariant algebras, and len instead of maj .

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$$w^{(3)} = 13.12.11223. \qquad \text{cdes}(w^{(3)}) = 3,$$

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$$W_{\alpha, \delta} := \text{words } w \text{ with content } \alpha \text{ and } \text{CDT}(w) = \delta.$$

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Theorem (Ahlbach-S.)

$(W_{\alpha,\delta}, \mathbb{Z}/n, W_{\alpha,\delta}^{\text{maj}}(q))$ refines the CSP triple $(W_{\alpha}, \mathbb{Z}/n, W_{\alpha}^{\text{maj}}(q))$.

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Remark

Completely different proof than RSW. Combinatorial and largely recursive. Involves Carlitz-style decomposition, (more or less new) notion of “modular periodicity,” a CSP extension lemma, a non-equivariant-but-fixed-point-preserving bijection, products of CSP’s on sets and multisets.

Refined CSP on Words

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Theorem (Ahlbach-S.)

Let $\alpha \vDash n$ be a strong composition with m parts, $\delta \vDash k$, $n_i := |w^{(i)}|$, $k_i := \text{cdes}(w^{(i)})$, $d := \text{gcd}(n, k)$. Then, modulo $q^n - 1$,

$$\begin{aligned} W_{\alpha, \delta}^{\text{maj}}(q) &\equiv \frac{d}{\alpha_1} [n/d]_{q^d} \prod_{\ell=2}^m q^{k_\ell \alpha_\ell} \binom{n_{\ell-1} - k_{\ell-1}}{\delta_\ell}_q \left(\binom{k_\ell}{\alpha_\ell - \delta_\ell} \right)_{q^{-1}} \\ &\equiv \frac{d}{\alpha_1} [n/d]_{q^d} q^\eta \prod_{\ell=2}^m \binom{n_{\ell-1} - k_{\ell-1}}{\delta_\ell}_q \left(\binom{k_\ell}{\alpha_\ell - \delta_\ell} \right)_q \end{aligned}$$

where $\eta := \binom{k}{2} + \sum_{\ell=2}^m \binom{\delta_\ell}{2} - \alpha_1$.

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Definition

Given $\lambda \vdash n - 1$, let $\lambda^\square \vdash n$ be the following “slightly skew partition”:

$$\lambda = \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \\ \hline \square & \square & \\ \hline \end{array} \Rightarrow \lambda^\square = \begin{array}{|c|c|c|c|} \hline & & & \square \\ \hline \square & \square & \square & \\ \hline \square & \square & & \\ \hline \square & \square & & \\ \hline \end{array}$$

Refined Sieving on Tableaux

Remark

Elizalde–Roichman (2017) defined a bijection

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- (i) $\text{cDes}(T) \cap [n-1] = \text{Des}(T)$
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Proof reduces to showing $[n]_q$ divides $\text{SYT}(\lambda^\square)^{\text{maj}}(q)$. Follows from

$$\begin{aligned} \text{SYT}(\lambda^\square)^{\text{maj}}(q) &= \binom{n}{n-1, 1}_q \text{SYT}(\lambda)^{\text{maj}}(q) \text{SYT}(\square)^{\text{maj}}(q) \\ &= [n]_q \text{SYT}(\lambda)^{\text{maj}}(q). \end{aligned}$$

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Remark

Showing $[n]_q \mid \text{SYT}(\lambda^\square; k)^{\text{maj}}(q)$ is significantly more involved. Uses an inner product formula of Adin–Reiner–Roichman (2017) for Elizalde–Roichman’s cyclic descent extensions, a “change of basis,” and the $W_{\alpha, \delta}^{\text{maj}}(q)$ product formula above.

Further work

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- ▶ In progress: further explore the CSP and Roichman et al's other cyclic descent extensions
- ▶ Give a representation-theoretic proof of $W_{\alpha,\delta}$ result

Thanks!

THANKS!